Subtropical Satisfiability

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SMT + non linear arithmetics

- High demand for non linear arithmetic reasoning capability
- Theory of real closed fields: decidable (QE: CAD, virtual substitution,...)
- Doubly exponential (existential fragment also high complexity)
- Complete decision procedure not always efficient enough
- Need for good heuristics

Our contribution

Simple heuristic to quickly discharge many proof obligations (or failing quickly)

- Based on subtropical method: quickly find positive solution for f = 0 where f has hundreds of thousand of monomials, with dozen variables, degrees around 10 in each variable
- Here: find real solution for $f_1 > 0 \land \cdots \land f_n > 0$

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- ▶ $-1 2x + x^3 > 0$, satisfiable: with sufficiently large x ▶ $2x - x^3 > 0$



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Subtropical method: towards the multivariate case

Polynomial $-2x_1^5+x_1^2x_2-3x_1^2-x_2^3+2x_2^2$ can be

- ▶ negative, e.g. $-2x_1^5$ dominates if x_1 large enough w.r.t. x_2
- ▶ positive, e.g. 2x₂² dominates if x₂ small enough (not zero) x₂ and an even smaller x₁.

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- reduce to univariate, setting all variables but one to 0
- consider monomial of highest/lowest total degree (if unique)
- ordering? lexicographic?

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Contribution

monotonic total preorders on the exponent vectors

Strictly max. monomials (w.r.t. these preorders) can dominate, for suitable (positive) values of variables.

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► $-2x_1^5$, $-3x_1^2$, $-x_2^3$ and $2x_2^2$ correspond to vertices



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-2x₁⁵, -3x₁², -x₂³ and 2x₂² correspond to vertices
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- ► $-2x_1^5$, $-3x_1^2$, $-x_2^3$ and $2x_2^2$ correspond to vertices
- These monomials can dominate for suitable values of variables
- Normal vector of separating plane provides witnesses



 $f = -2x_1^5 + x_1^2x_2 - 3x_1^2 - x_2^3 + 2x_2^2$

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Subtropical method: from preorders to QF_LRA SMT



monotonic total preorders correspond to normal vectors

- $(x_1, x_2) \preceq (x'_1, x'_2)$ iff $-3x_1 2x_2 \le -3x'_1 2x'_2$
- $(5,0) \prec (2,1) \prec (2,0) \approx (0,3) \prec (0,2)$

To QF_LRA SMT?

• $S^+ = \{(2,1), (0,2)\}, S^- = \{(5,0), (2,0), (0,3)\}$

$$f>0 \longleftarrow \bigwedge_{(p_1,p_2)\in \mathcal{S}^-} p_1n_1 + p_2n_2 + c < 0 \land \bigvee_{(p_1,p_2)\in \mathcal{S}^+} p_1n_1 + p_2n_2 + c > 0$$

linear constraints on real variables, n₁, n₂, c

One polynomial:

$$f_1 = -2x_1^5 + x_1^2x_2 -3x_1^2 - x_2^3 + 2x_2^2$$

One polynomial:

build the Newton polytope



One polynomial:

- build the Newton polytope
- find a suitable vertex



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$$\begin{array}{ll} f_1 = -2x_1^5 + x_1^2 x_2 & f_2 = 1 - x_1 x_2 & f_3 = x_1 x_2 \\ -3x_1^2 - x_2^3 + 2x_2^2 & -5x_1 - 6x_2 & -x_1^5 x_2^2 + x_1 x_2^4 \end{array}$$

Several polynomials:

One polynomial:

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- find a suitable vertex
- normal vector to separating plane provides witness

Several polynomials:

build the Newton polytopes



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common normal vector ensures existence of global solution

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n polynomial constraints? Conjunction of n QF_LRA problems sharing only variables to describe normal vector

• Up to now: $\bigwedge_i f_i > 0$ with all $\bigwedge_i x_i > 0$



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- This can be encoded into the QF_LRA SMT problem; no need to check 2ⁿ formulas



Experimental results

- STROPSAT integrated in veriT (not the SMT-COMP version)
- Tested on SMT-LIB/QF_NRA on suitable problems, i.e. 4917/11601 files: 3265 sat, 106 unknown, 1546 unsat
- CVC4 used to handle linear solving
- 2500s timeout, 20GB

On 1546 unsat-labeled formulas: 200 unsat by LRA, cumulative time to fail on the 1346 others: 18.45s, max 0.1s Shows satisfiability for 2403 problems, including 15 "unknown" problems (and 9 where Z3 fails)

Experimental results



When STROPSAT does not fail

- time comparable to Z3
- sometimes succeeds alone
- if timeouts, Z3 too

STROPSAT is quick to fail

Conclusion

- ► A heuristic, providing quick solutions, or failing quickly
- Good results for many SMT benchmarks
- Not sensitive to the number of variables; actually, gets "better" when the number of variables grows
- Investigate its use in context where getting models is paramount, i.e. testing phase of raSAT loop
- What can we do along these lines to help complete decision procedures?
- Better understand when the method works

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